

# Chapter 5 Finite Difference Methods

# References

1. Chapters 5 and 9, Brandimarte
2. Section 17.8, Hull
3. Chapter 7, “Numerical analysis”, Burden and Faires

# Outline

- Finite difference (FD) approximation to the derivatives
- Explicit FD method
- Numerical issues
- Implicit FD method
- Crank-Nicolson method
- Dealing with American options
- Further comments

# Chapter 5 Finite Difference Methods

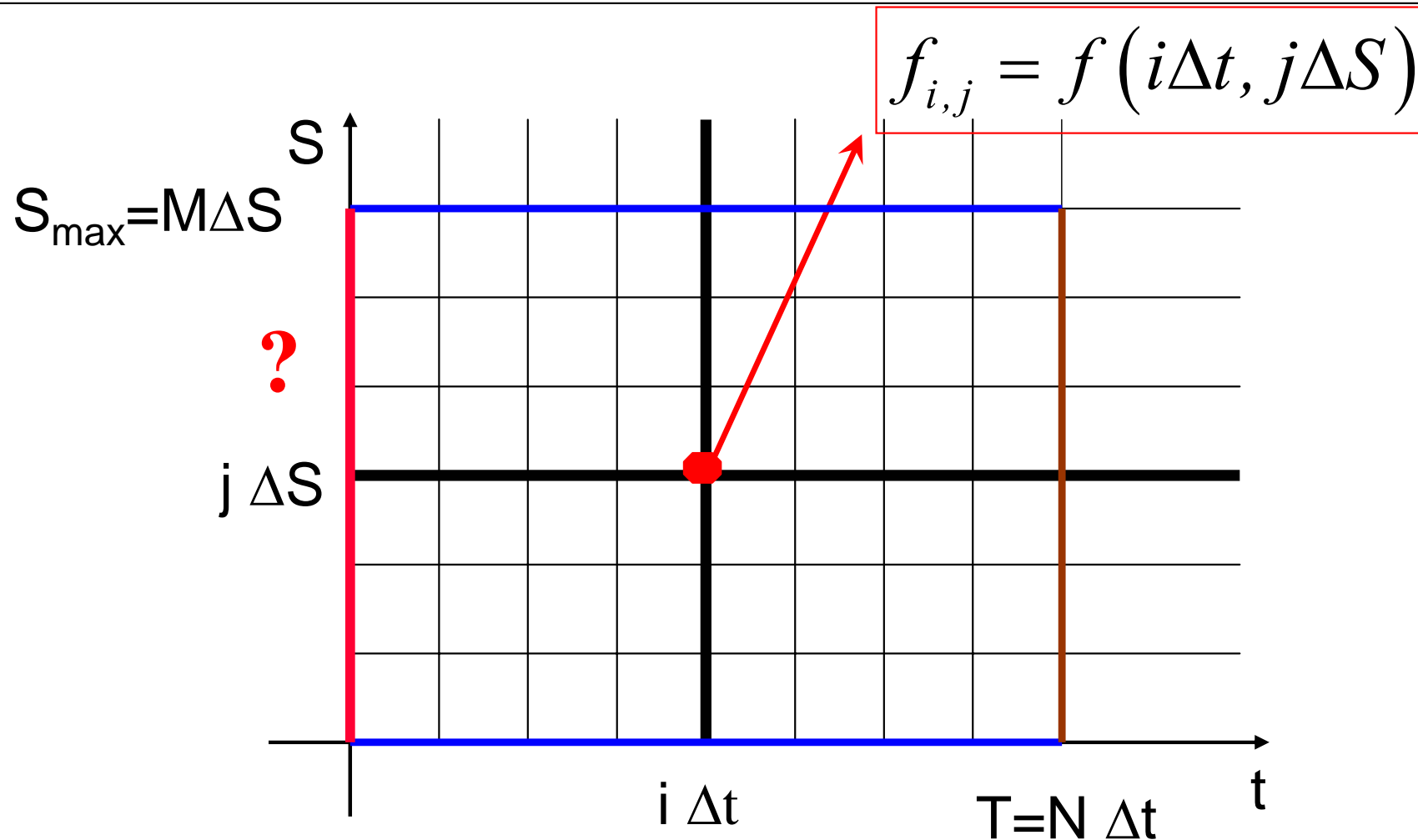
## 5.1 Finite difference approximations

# Finite-difference mesh

- Aim to approximate the values of the continuous function  $f(t, S)$  on a set of discrete points in  $(t, S)$  plane
- Divide the  $S$ -axis into equally spaced nodes at distance  $\Delta S$  apart, and, the  $t$ -axis into equally spaced nodes a distance  $\Delta t$  apart
- $(t, S)$  plane becomes a mesh with mesh points on  $(i \Delta t, j \Delta S)$
- We are interested in the values of  $f(t, S)$  at mesh points  $(i \Delta t, j \Delta S)$ , denoted as

$$f_{i,j} = f(i\Delta t, j\Delta S)$$

# The mesh for finite-difference approximation



# Black-Scholes Equation for a European option with value $V(S,t)$

$$\frac{\partial V}{\partial t} + \frac{1}{2} \sigma^2 S^2 \frac{\partial^2 V}{\partial S^2} + rS \frac{\partial V}{\partial S} - rV = 0 \quad (5.1)$$

where  $0 < S < +\infty$  and  $0 \leq t < T$

with proper final and boundary conditions

## Notes:

This is a second-order **hyperbolic, elliptic, or parabolic, forward or backward** partial differential equation

Its solution is sufficiently well behaved ,i.e. well-posed

# Finite difference approximations

The basic idea of FDM is to replace the partial derivatives by approximations obtained by Taylor expansions near the point of interests

For example,

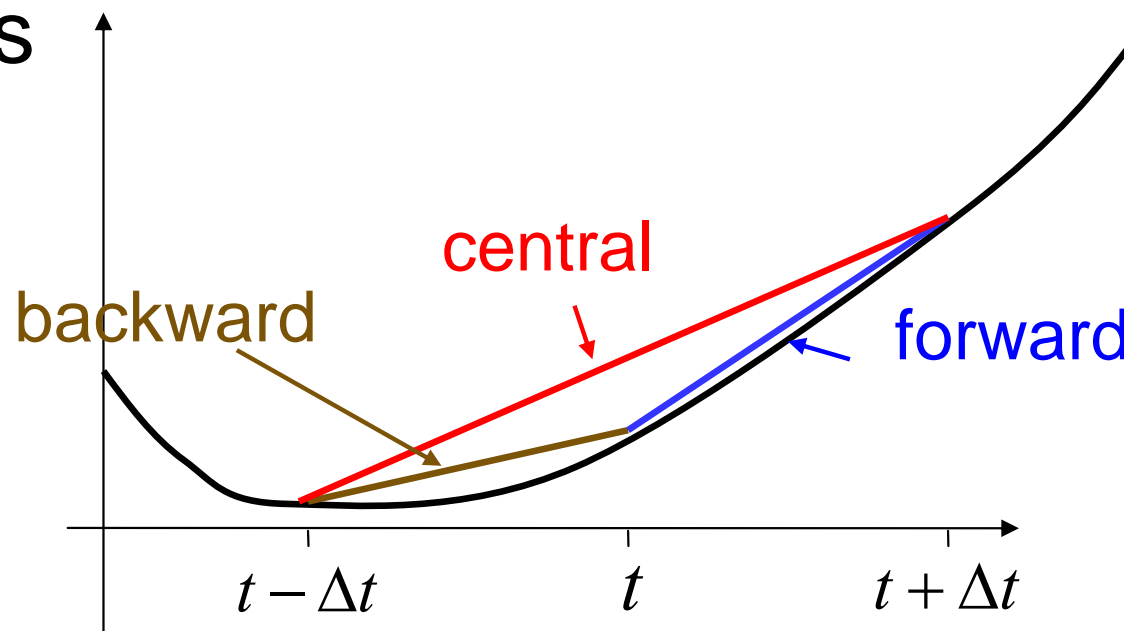
$$\frac{\partial f(S, t)}{\partial t} = \lim_{\Delta t \rightarrow 0} \frac{f(S, t + \Delta t) - f(S, t)}{\Delta t} \approx \frac{f(S, t + \Delta t) - f(S, t)}{\Delta t}$$

for small  $\Delta t$ , using Taylor expansion at point  $(S, t)$

$$f(S, t + \Delta t) = f(S, t) + \frac{\partial f(S, t)}{\partial t} \Delta t + O((\Delta t)^2)$$



# Forward-, Backward-, and Central-difference approximation to 1<sup>st</sup> order derivatives



Forward: 
$$\frac{\partial f(t, S)}{\partial t} \approx \frac{f(t + \Delta t, S) - f(t, S)}{\Delta t} + O(\Delta t)$$

Backward: 
$$\frac{\partial f(t, S)}{\partial t} \approx \frac{f(t, S) - f(t - \Delta t, S)}{\Delta t} + O(\Delta t)$$

Central: 
$$\frac{\partial f(t, S)}{\partial t} \approx \frac{f(t + \Delta t, S) - f(t - \Delta t, S)}{2\Delta t} + O((\Delta t)^2)$$

# Symmetric Central-difference approximations to 2<sup>nd</sup> order derivatives

$$\frac{\partial^2 f(t, S)}{\partial S^2} \approx \frac{f(t, S + \Delta S) - 2f(t, S) + f(t, S - \Delta S)}{(\Delta S)^2} + O((\Delta S)^2)$$

Use Taylor's expansions for  $f(t, S + \Delta S)$  and  $f(t, S - \Delta S)$   
around point  $(t, S)$ :

$$f(t, S + \Delta S) = ?$$

+

$$f(t, S - \Delta S) = ?$$

# Finite difference approximations

Forward Difference:  $\frac{\partial f}{\partial t} \approx \frac{f_{i+1,j} - f_{i,j}}{\Delta t}, \quad \frac{\partial f}{\partial S} \approx \frac{f_{i,j+1} - f_{i,j}}{\Delta S}$

Backward Difference:  $\frac{\partial f}{\partial t} \approx \frac{f_{i,j} - f_{i-1,j}}{\Delta t}, \quad \frac{\partial f}{\partial S} \approx \frac{f_{i,j} - f_{i,j-1}}{\Delta S}$

Central Difference:  $\frac{\partial f}{\partial t} \approx \frac{f_{i+1,j} - f_{i-1,j}}{2\Delta t}, \quad \frac{\partial f}{\partial S} \approx \frac{f_{i,j+1} - f_{i,j-1}}{2\Delta S}$

As to the second derivative, we have:

$$\begin{aligned} \frac{\partial^2 f}{\partial S^2} &\approx \left( \frac{f_{i,j+1} - f_{i,j}}{\Delta S} - \frac{f_{i,j} - f_{i,j-1}}{\Delta S} \right) / \Delta S \\ &= \frac{f_{i,j+1} - 2f_{i,j} + f_{i,j-1}}{(\Delta S)^2} \end{aligned}$$

# Finite difference approximations

- Depending on which combination of schemes we use in discretizing the equation, we will have explicit, implicit, or Crank-Nicolson methods
- We also need to discretize the boundary and final conditions accordingly. For example, for European Call,

Final Condition:

$$f_{N,j} = \max(j\Delta S - K, 0), \quad \text{for } j = 0, 1, \dots, M$$

Boundary Conditions:

$$\begin{cases} f_{i,0} = 0 \\ f_{i,M} = S_{\max} - Ke^{-r(N-i)\Delta t} \end{cases}, \quad \text{for } i = 0, 1, \dots, N$$

where  $S_{\max} = M\Delta S$ .

# Chapter 5 Finite Difference Methods

## 5.2.1 Explicit Finite-Difference Method

# Explicit Finite Difference Methods

In  $\frac{\partial f}{\partial t} + rS \frac{\partial f}{\partial S} + \frac{1}{2} \sigma^2 S^2 \frac{\partial^2 f}{\partial S^2} = rf$ , at point  $(i\Delta t, j\Delta S)$ , set

backward difference:  $\frac{\partial f}{\partial t} \approx \frac{f_{i,j} - f_{i-1,j}}{\Delta t}$

central difference:  $\frac{\partial f}{\partial S} \approx \frac{f_{i,j+1} - f_{i,j-1}}{2\Delta S}$ ,

and

$$\frac{\partial^2 f}{\partial S^2} \approx \frac{f_{i,j+1} + f_{i,j-1} - 2f_{i,j}}{\Delta S^2}, \quad rf = rf_{i,j}, \quad S = j\Delta S$$

## Explicit Finite Difference Methods

Rewriting the equation, we get an explicit scheme:

$$f_{i-1,j} = a_j^* f_{i,j-1} + b_j^* f_{i,j} + c_j^* f_{i,j+1} \quad (5.2)$$

where

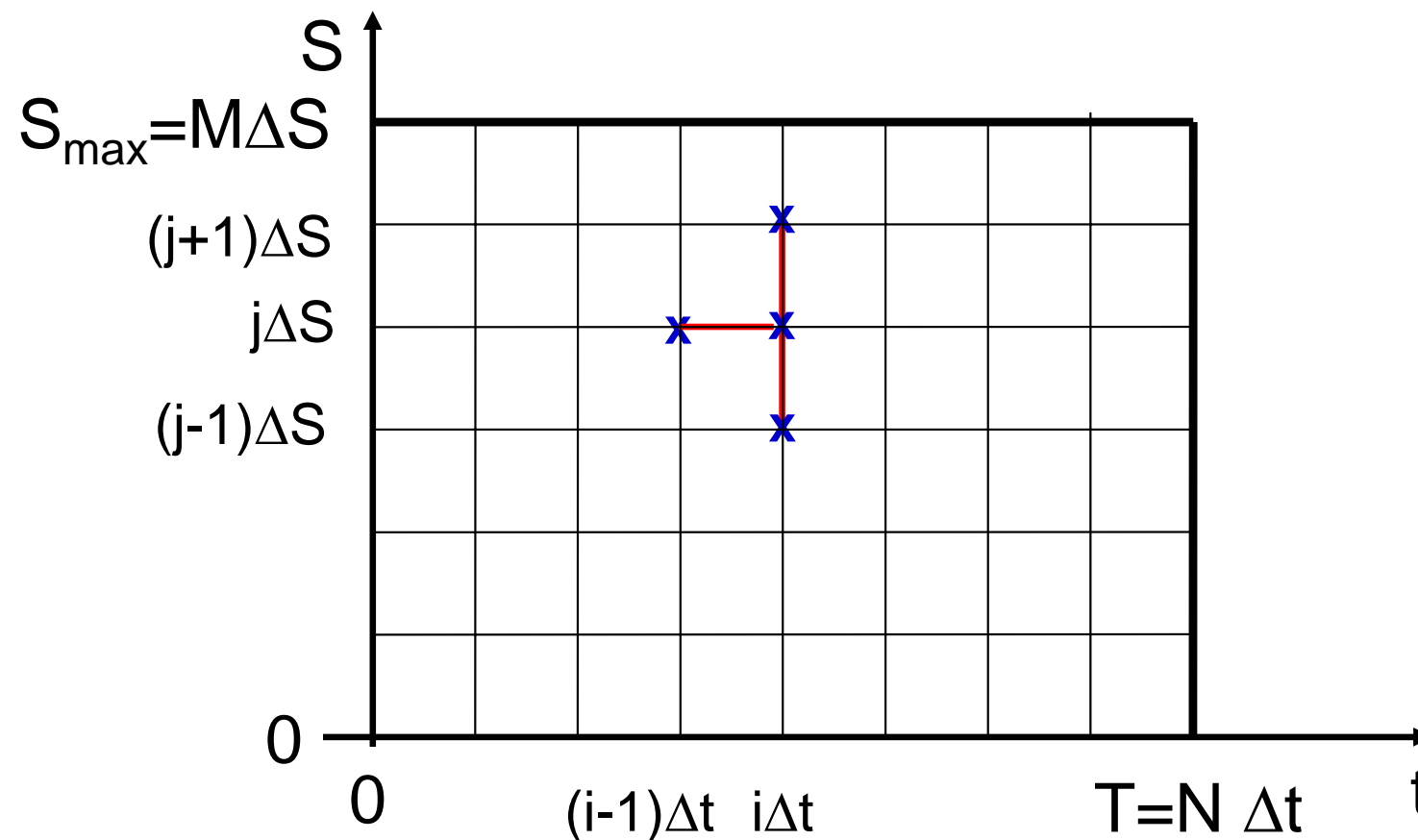
$$a_j^* = \frac{1}{2} \Delta t (\sigma^2 j^2 - rj)$$

$$b_j^* = 1 - \Delta t (\sigma^2 j^2 + r)$$

$$c_j^* = \frac{1}{2} \Delta t (\sigma^2 j^2 + rj)$$

for  $i = N - 1, N - 2, \dots, 1, 0$  and  $j = 1, 2, \dots, M - 1$ .

# Numerical Computation Dependency





# Implementation

1. Starting with the final values  $f_{N,j}$ , we apply (5.2) to solve  $f_{N-1,j}$  for  $1 \leq j \leq M-1$ . We use the boundary condition to determine  $f_{N-1,0}$  and  $f_{N-1,M}$ .
2. Repeat the process to determine  $f_{N-2,j}$  and so on

# Example

We compare explicit finite difference solution for a European put with the exact Black-Scholes formula, where  $T = 5/12$  yr,  $S_0 = \$50$ ,  $K = \$50$ ,  $\sigma = 30\%$ ,  $r = 10\%$ .

Black-Scholes Price: \$2.8446

EFD Method with  $S_{\max} = \$100$ ,  $\Delta S = 2$ ,  $\Delta t = 5/1200$ : \$2.8288

EFD Method with  $S_{\max} = \$100$ ,  $\Delta S = 1$ ,  $\Delta t = 5/4800$ : \$2.8406

## Example (Stability)

We compare explicit finite difference solution for a European put with the exact Black-Scholes formula, where  $T = 5/12$  yr,  $S_0 = \$50$ ,  $K = \$50$ ,  $\sigma = 30\%$ ,  $r = 10\%$ .

Black-Scholes Price: \$2.8446

EFD Method with  $S_{\max} = \$100$ ,  $\Delta S = 2$ ,  $\Delta t = 5/1200$ : \$2.8288

EFD Method with  $S_{\max} = \$100$ ,  $\Delta S = 1.5$ ,  $\Delta t = 5/1200$ : \$3.1414

EFD Method with  $S_{\max} = \$100$ ,  $\Delta S = 1$ ,  $\Delta t = 5/1200$ :  $-\$2.8271E22$

# Chapter 5 Finite Difference Methods

## 5.2.2 Numerical Stability

# Numerical Accuracy

- The problem itself
- The discretization scheme used
- The numerical algorithm used

# Conditioning Issue

Suppose we have mathematically posed problem:

$$y = f(x)$$

where  $y$  is to be evaluated given an input  $x$ .

Let  $x^* = x + \delta x$  for small change  $\delta x$ .

If  $f(x^*)$  is near  $f(x)$ , then we call the problem is *well-conditioned*. Otherwise, it is ill-posed/ill-conditioned.

# Conditioning Issue

- Conditioning issue is related to the problem itself, not to the specific numerical algorithm; Stability issue is related to the numerical algorithm
- One can not expect a good numerical algorithm to solve an ill-conditioned problem any more accurately than the data warrant
- But a bad numerical algorithm can produce poor solutions even to well-conditioned problems

# Conditional Issue

The concept "near" can be measured by further information about the particular problem:

$$\frac{\|f(x) - f(x^*)\|}{\|f(x)\|} \leq C \frac{\|\delta x\|}{\|x\|} \quad (f(x) \neq 0)$$

where  $C$  is called *condition number* of this problem. If  $C$  is large, the problem is ill-conditioned.



# Floating Point Number & Error

Let  $x$  be any real number.

Infinite decimal expansion :  $x = \pm .x_1x_2 \cdots x_d \cdots 10^e$

Truncated floating point number :  $x \approx fl(x) = \pm .x_1x_2 \cdots x_d 10^e$

where  $x_1 \neq 0, 0 \leq x_i \leq 9$ ,

$d$  : an integer, precision of the floating point system

$e$  : an bounded integer

**Floating point or roundoff error :  $fl(x) - x$**

# Error Propagation

When additional calculations are done, there is an accumulation of these floating point errors.

**Example :** Let  $x = -0.6667$  and  $fl(x) = -0.667 \cdot 10^0$  where  $d = 3$ .

Floating point error :  $fl(x) - x = -0.0003$

Error propagation :  $fl(x)^2 - x^2 = 0.00040011$

# Numerical Stability or Instability

Stability ensures if the error between the numerical solution and the exact solution remains bounded as numerical computation progresses.

That is,  $f(x^*)$  (the solution of a slightly perturbed problem) is near  $f^*(x)$  (the computed solution).

Stability concerns about the behavior of  $|f_{i,j} - f(i\Delta t, j\Delta S)|$  as numerical computation progresses for fixed discretization steps  $\Delta t$  and  $\Delta S$ .

# Convergence issue

Convergence of the numerical algorithm concerns about the behavior of  $|f_{i,j} - f(i\Delta t, j\Delta S)|$  as  $\Delta t, \Delta S \rightarrow 0$  for fixed values  $(i\Delta t, j\Delta S)$ .

For well-posed linear initial value problem,

Stability  $\Leftrightarrow$  Convergence

(Lax's equivalence theorem, Richtmyer and Morton, "Difference Methods for Initial Value Problems" (2nd) 1967)

# Numerical Accuracy

- These factors contribute the accuracy of a numerical solution. We can find a “good” estimate if our problem is well-conditioned and the algorithm is stable

$$\left. \begin{array}{l} \text{Stable: } f^*(x) \approx f(x^*) \\ \text{Well-conditioned: } f(x^*) \approx f(x) \end{array} \right\} f^*(x) \approx f(x)$$

# Chapter 5 Finite Difference Methods

## 5.2.3 Financial Interpretation of Numerical Instability

## Financial Interpretation of instability (Hall, page 423-4)

If  $\partial f / \partial S$  and  $\partial^2 f / \partial S^2$  are assumed to be the same at  $(i+1, j)$  as they are at  $(i, j)$ , we obtain equations of the form:

$$f_{i,j} = \hat{a}_j f_{i+1,j-1} + \hat{b}_j f_{i+1,j} + \hat{c}_j f_{i+1,j+1} \quad (5.3)$$

where

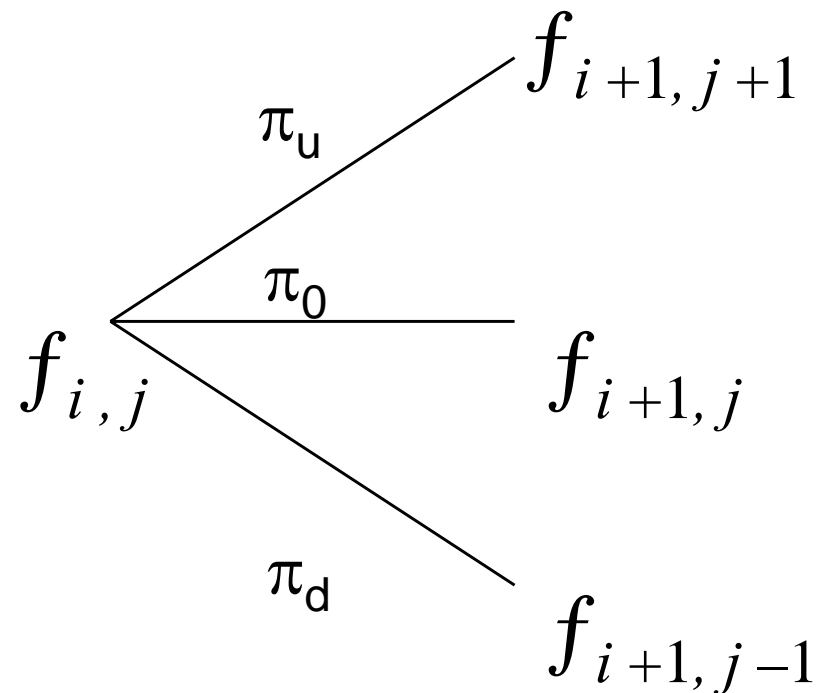
$$\hat{a}_j = \frac{1}{1+r\Delta t} \left( \frac{1}{2} \Delta t \sigma^2 j^2 - \frac{1}{2} \Delta t r j \right) = \frac{1}{1+r\Delta t} \pi_d$$

$$\hat{b}_j = \frac{1}{1+r\Delta t} (1 - \sigma^2 j^2 \Delta t) = \frac{1}{1+r\Delta t} \pi_0$$

$$\hat{c}_j = \frac{1}{1+r\Delta t} \left( \frac{1}{2} \Delta t \sigma^2 j^2 + \frac{1}{2} \Delta t r j \right) = \frac{1}{1+r\Delta t} \pi_u$$

for  $i = N-1, N-2, \dots, 1, 0$  and  $j = 1, 2, \dots, M-1$ .

# Explicit Finite Difference Methods



These coefficients can be interpreted as probabilities times a discount factor. If one of these probability  $< 0$ , instability occurs.



# Explicit Finite Difference Method as Trinomial Tree

Check if the mean and variance of the

Expected value of the increase in asset price during  $\Delta t$ :

$$E[\Delta] = -\Delta S \pi_d + 0\pi_0 + \Delta S \pi_u = rj\Delta S \Delta t = rS\Delta t$$

Variance of the increment:

$$E[\Delta^2] = -(\Delta S)^2 \pi_d + 0\pi_0 + (\Delta S)^2 \pi_u = \sigma^2 j^2 (\Delta S)^2 \Delta t = \sigma^2 S^2 \Delta t$$

$$\text{Var}[\Delta] = E[\Delta^2] - E^2[\Delta] = \sigma^2 S^2 \Delta t - r^2 S^2 (\Delta t)^2 \approx \sigma^2 S^2 \Delta t$$

which is coherent with geometric Brownian motion in a risk-neutral world

# Change of Variable

Define  $Z = \ln S$ . The B-S equation becomes

$$\frac{\partial f}{\partial t} + \left( r - \frac{\sigma^2}{2} \right) \frac{\partial f}{\partial Z} + \frac{\sigma^2}{2} \frac{\partial^2 f}{\partial Z^2} = rf$$

The corresponding difference equation is

$$\frac{f_{i+1,j} - f_{i,j}}{\Delta t} + \left( r - \frac{\sigma^2}{2} \right) \frac{f_{i+1,j+1} - f_{i+1,j-1}}{2\Delta Z} + \frac{\sigma^2}{2} \frac{f_{i+1,j+1} - 2f_{i+1,j} + f_{i+1,j-1}}{\Delta Z^2} = rf_{i,j}$$

or

$$f_{i,j} = \alpha_j^* f_{i+1,j-1} + \beta_j^* f_{i+1,j} + \gamma_j^* f_{i+1,j+1} \quad (5.4)$$

# Change of Variable

where

$$\alpha_j^* = \frac{1}{1+r\Delta t} \left[ - \left( r - \frac{\sigma^2}{2} \right) \frac{\Delta t}{2\Delta Z} + \frac{\sigma^2}{2} \frac{\Delta t}{\Delta Z^2} \right]$$

$$\beta_j^* = \frac{1}{1+r\Delta t} \left( 1 - \frac{\Delta t}{\Delta Z^2} \sigma^2 \right)$$

$$\gamma_j^* = \frac{1}{1+r\Delta t} \left[ \left( r - \frac{\sigma^2}{2} \right) \frac{\Delta t}{2\Delta Z} + \frac{\sigma^2}{2} \frac{\Delta t}{\Delta Z^2} \right]$$

It can be shown that it is numerically most efficient if  $\Delta Z = \sigma\sqrt{3\Delta t}$ .

# Reduced to Heat Equation

Get rid of the varying coefficients  $S$  and  $S^2$  by using change of variables:

$$S = Ee^x, t = T - 2\tau/\sigma^2, V(S, t) = Ee^{-\frac{1}{2}(k-1)x - \frac{1}{4}(k+1)^2\tau} u(x, \tau)$$

$$k = r / \left( \frac{1}{2} \sigma^2 \right)$$

Equation (5.1) becomes heat equation (5.5):

$$\frac{\partial u}{\partial \tau} = \frac{\partial^2 u}{\partial x^2} \quad (5.5)$$

for  $-\infty < x < +\infty$  and  $\tau > 0$

# Explicit Finite Difference Method

With  $u_n^m = u(n\delta x, m\delta\tau)$ , this involves solving a system of finite difference equations of the form :

$$\frac{u_n^{m+1} - u_n^m}{\delta\tau} + O(\delta\tau) = \frac{u_{n+1}^m - 2u_n^m + u_{n-1}^m}{(\delta x)^2} + O((\delta x)^2)$$

Ignoring terms of  $O(\delta\tau)$  and  $O((\delta x)^2)$ , we can approximate this by

$$u_n^{m+1} = \alpha u_{n+1}^m + (1 - 2\alpha)u_n^m + \alpha u_{n-1}^m$$

where  $\alpha = \frac{\delta\tau}{(\delta x)^2}$ , for  $-N^- \leq n \leq N^+$  and  $m = 0, 1, \dots, M$   $\left( = \frac{\frac{1}{2}\sigma^2 T}{\delta\tau} \right)$

# Stability and Convergence

(P. Wilmott, et al, Option Pricing)

## Stability:

The solution of Eqn (5.5) is

$$\text{i) Stable if } 0 < \alpha = \frac{\delta\tau}{(\delta x)^2} \leq \frac{1}{2}; \quad \text{ii) Unstable if } \alpha > \frac{1}{2}$$

## Convergence:

If  $0 < \alpha \leq \frac{1}{2}$ , then the explicit finite-difference approximation

converges to the exact solution as  $\delta\tau, \delta x \rightarrow 0$

(in the sense that  $u_n^m \rightarrow u(n\delta x, m\delta\tau)$  as  $\delta\tau, \delta x \rightarrow 0$ )

**Rate of Convergence is  $O(\delta\tau)$**

# Chapter 5 Finite Difference Methods

## 5.3.1 Implicit Finite-Difference Method

# Implicit Finite Difference Methods

In  $\frac{\partial f}{\partial t} + rS \frac{\partial f}{\partial S} + \frac{1}{2} \sigma^2 S^2 \frac{\partial^2 f}{\partial S^2} = r f$ , we use

**forward difference:**  $\frac{\partial f}{\partial t} \approx \frac{f_{i+1,j} - f_{i,j}}{\Delta t}$

**central difference:**  $\frac{\partial f}{\partial S} \approx \frac{f_{i,j+1} - f_{i,j-1}}{2\Delta S}$ ,

and

$$\frac{\partial^2 f}{\partial S^2} \approx \frac{f_{i,j+1} + f_{i,j-1} - 2f_{i,j}}{\Delta S^2}, \quad rf = rf_{i,j}$$



# Implicit Finite Difference Methods

Rewriting the equation, we get an implicit scheme:

$$a_j f_{i,j-1} + b_j f_{i,j} + c_j f_{i,j+1} = f_{i+1,j} \quad (5.6)$$

where

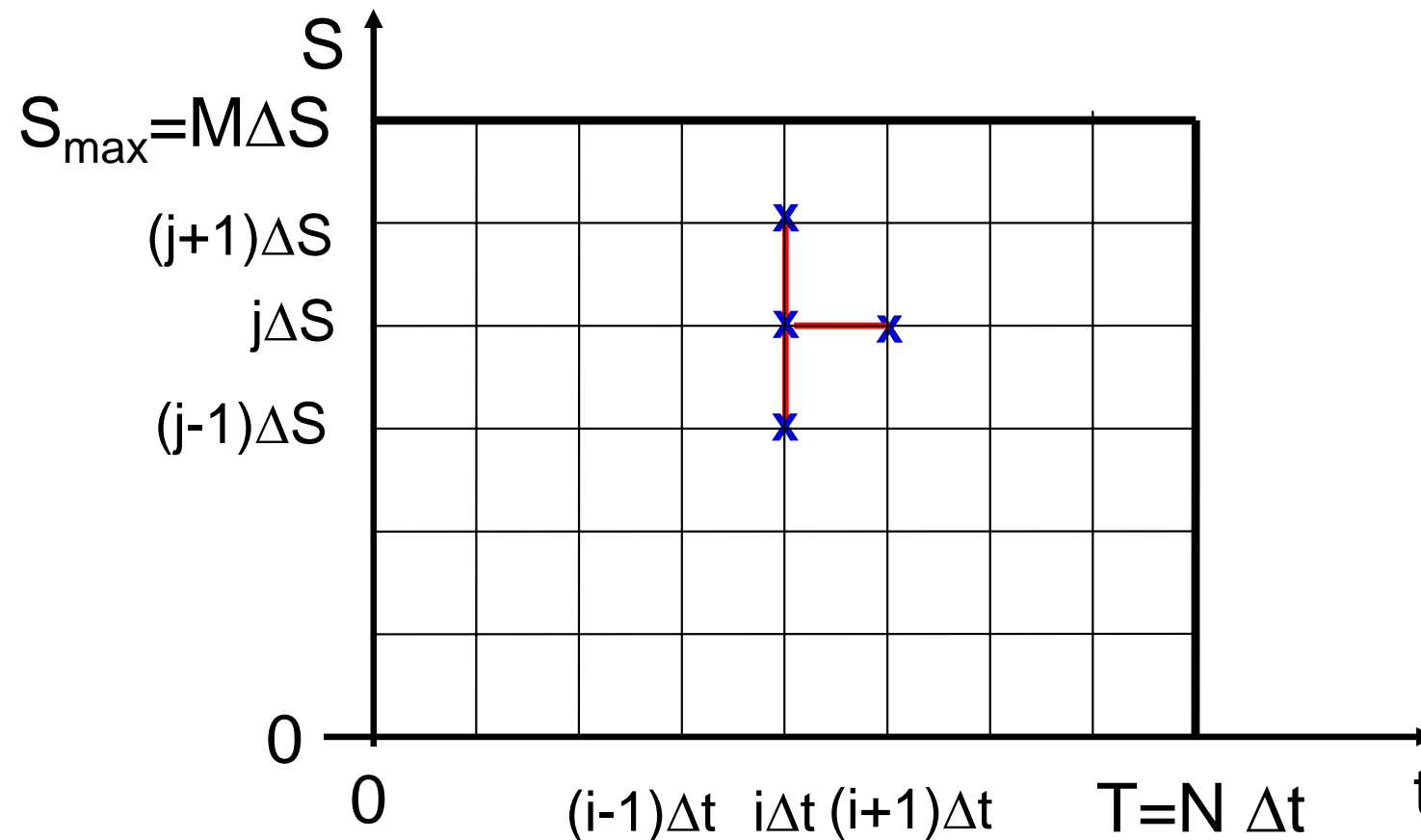
$$a_j = \frac{1}{2} \Delta t \left( -\sigma^2 j^2 + rj \right)$$

$$b_j = 1 + \Delta t \left( \sigma^2 j^2 + r \right)$$

$$c_j = -\frac{1}{2} \Delta t \left( \sigma^2 j^2 + rj \right)$$

for  $i = N - 1, N - 2, \dots, 1, 0$  and  $j = 1, 2, \dots, M - 1$ .

# Numerical Computation Dependency



# Implementation

Equation (5.6) can be rewritten in matrix form:

$$\mathbf{C}\mathbf{f}_i = \mathbf{f}_{i+1} + \mathbf{b}_i \quad (5.7)$$

where  $\mathbf{f}_i$  and  $\mathbf{b}_i$  are  $(M-1)$  dimensional vectors

$$\mathbf{f}_i = [f_{i,1}, f_{i,2}, f_{i,3} \cdots, f_{i,M-1}]^T, \mathbf{b}_i = [-a_1 f_{i,0}, 0, 0, \cdots, 0, -c_{M-1} f_{i,M}]^T$$

and  $\mathbf{C}$  is  $(M-1) \times (M-1)$  symmetric matrices

$$\mathbf{C} = \begin{bmatrix} b_1 & c_1 & 0 & \cdots & 0 \\ a_2 & b_2 & c_2 & \ddots & 0 \\ 0 & a_3 & b_3 & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & c_{M-2} \\ 0 & \cdots & 0 & a_{M-1} & b_{M-1} \end{bmatrix}$$

# Implementation

1. Starting with the final values  $f_{N,j}$ , we need to solve a linear system (5.7) to obtain  $f_{N-1,j}$  for  $1 \leq j \leq M-1$  using LU factorization or iterative methods. We use the boundary condition to determine  $f_{N-1,0}$  and  $f_{N-1,M}$ .
2. Repeat the process to determine  $f_{N-2,j}$  and so on

# Example

We compare implicit finite difference solution for a European put with the exact Black-Scholes formula, where  $T = 5/12$  yr,  $S_0 = \$50$ ,  $K = \$50$ ,  $\sigma = 30\%$ ,  $r = 10\%$ .

Black-Scholes Price: \$2.8446

IFD Method with  $S_{\max} = \$100$ ,  $\Delta S = 2$ ,  $\Delta t = 5/1200$ : \$2.8194

IFD Method with  $S_{\max} = \$100$ ,  $\Delta S = 1$ ,  $\Delta t = 5/4800$ : \$2.8383

## Example (Stability)

We compare implicit finite difference solution for a European put with the exact Black-Scholes formula, where  $T = 5/12$  yr,  $S_0 = \$50$ ,  $K = \$50$ ,  $\sigma = 30\%$ ,  $r = 10\%$ .

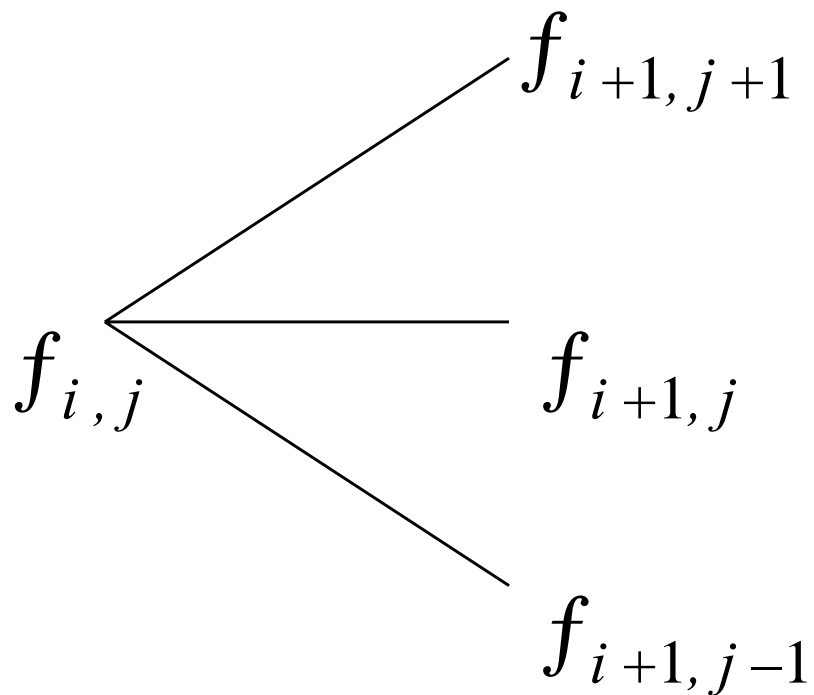
Black-Scholes Price: \$2.8846

IFD Method with  $S_{\max} = \$100$ ,  $\Delta S = 2$ ,  $\Delta t = 5/1200$ : \$2.8288

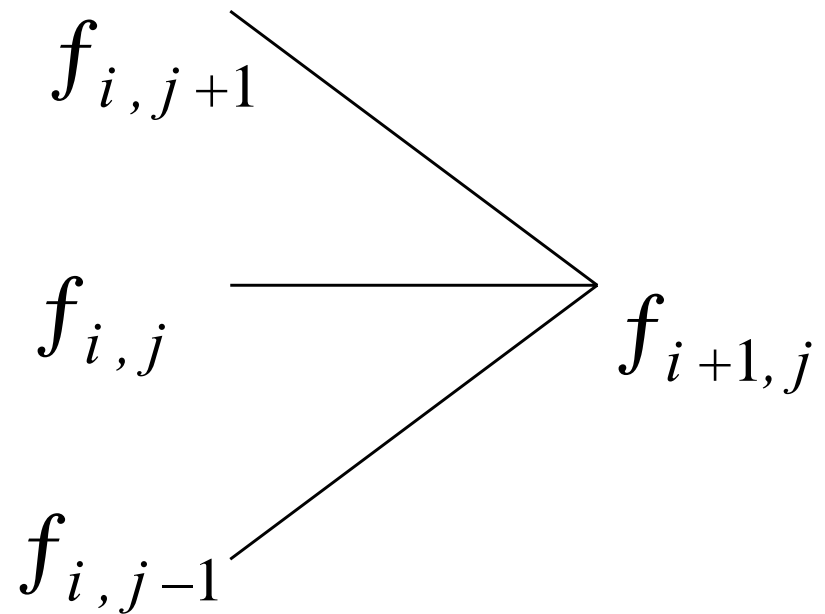
IFD Method with  $S_{\max} = \$100$ ,  $\Delta S = 1.5$ ,  $\Delta t = 5/1200$ : \$3.1325

IFD Method with  $S_{\max} = \$100$ ,  $\Delta S = 1$ ,  $\Delta t = 5/1200$ : \$2.8348

# Implicit vs Explicit Finite Difference Methods



Explicit Method



Implicit Method  
(always stable)

# Implicit vs Explicit Finite Difference Method

- The explicit finite difference method is equivalent to the trinomial tree approach:
  - Truncation error:  $O(\Delta t)$
  - Stability: not always
- The implicit finite difference method is equivalent to a multinomial tree approach:
  - Truncation error:  $O(\Delta t)$
  - Stability: always



# Other Points on Finite Difference Methods

- It is better to have  $\ln S$  rather than  $S$  as the underlying variable in general
- Improvements over the basic implicit and explicit methods:
  - Crank-Nicolson method, average of explicit and implicit FD methods, trying to achieve
    - Truncation error:  $O((\Delta t)^2)$
    - Stability: always

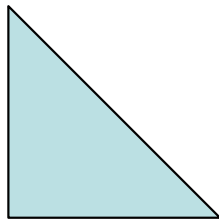
# Chapter 5 Finite Difference Methods

## 5.3.2 Solving a linear system using direct methods

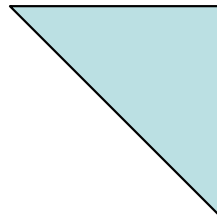
# Solve $Ax=b$

$$A x=b$$

## Various shapes of matrix $A$



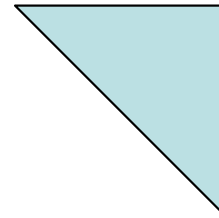
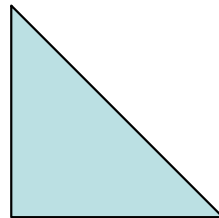
Lower triangular



Upper triangular



General



## 5.3.2.A Triangular Solvers

## Example: 3 x 3 upper triangular system

$$\begin{bmatrix} 4 & 6 & 1 \\ 0 & 1 & 1 \\ 0 & 0 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 100 \\ 10 \\ 20 \end{bmatrix}$$

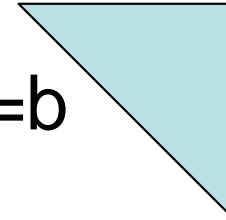
$$\Rightarrow x_3 = 20 / 4 = 5$$

$$\Rightarrow x_2 = 10 - x_3 = 5$$

$$\Rightarrow 4x_1 = 100 - x_3 - 6 * x_2 = 65$$

$$\therefore x_1 = 65/4$$

# Solve an upper triangular system $Ax=b$



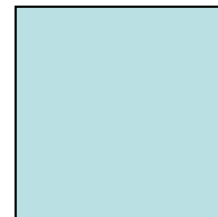
$$\begin{pmatrix} 0 & \cdots & 0 & a_{ii} & \cdots & a_{in} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} = a_{ii}x_i + \sum_{j>i} a_{ij}x_j = b_i, i = 1, \dots, n$$

$$\Rightarrow x_n = b_n / a_{nn}$$

$$x_i = \left( b_i - \sum_{j>i} a_{ij}x_j \right) / a_{ii}, i = 1, \dots, n-1$$

# Implementation

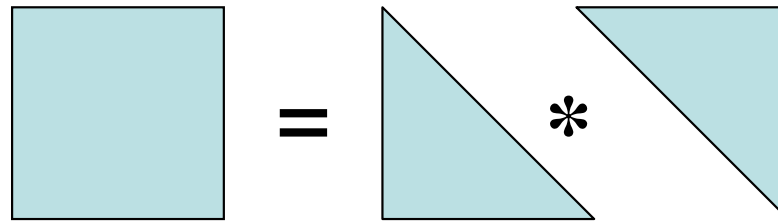
```
Function x = UpperTriSolve(A, b)
    n = length(b);
    x(n) = b(n)/A(n,n);
    for i = n-1:-1:1
        sum = b(i);
        for j = i+1:n
            sum = sum - A(i,j)*x(j);
        end
        x(i) = sum/A(i,i);
    end
end
```



## 5.3.2.B Gauss Elimination



To solve  $Ax=b$  for general form  $A$



To solve  $Ax = b$ :

Suppose  $A = LU$ . Then

$$Ax = LUx = b$$

$\Downarrow$

$z$

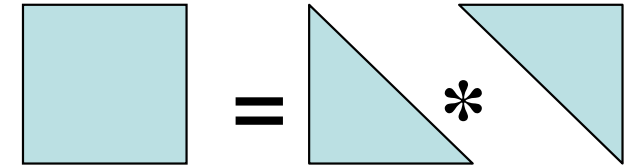


Solve two triangular systems:

1) solve  $z$  from  $Lz = b$

2) solve  $x$  from  $Ux = z$

# Gauss Elimination



**Goal:** Make A an upper triangular matrix through using fundamental row operations

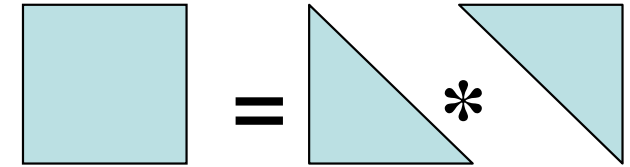
For example, to zero the elements in the lower triangular part of A

1) Zero  $a_{21}$

$$\begin{bmatrix} 1 & 0 & 0 \\ -\frac{a_{21}}{a_{11}} & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix} = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ 0 & a_{22} - \frac{a_{21}a_{12}}{a_{11}} & a_{23} - \frac{a_{21}a_{13}}{a_{11}} \\ a_{31} & a_{32} & a_{33} \end{bmatrix}$$

$\Downarrow$   
 $\mathbf{E}_{21}\mathbf{A}$

# Gauss Elimination



2) Zero  $a_{31}$

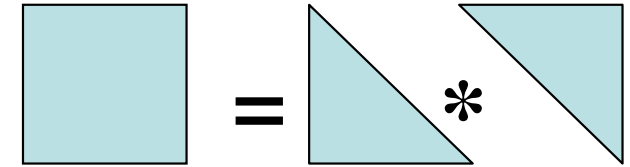
$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -\frac{a_{31}}{a_{11}} & 0 & 1 \end{bmatrix} \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ 0 & a_{22} - \frac{a_{21}a_{12}}{a_{11}} & a_{23} - \frac{a_{21}a_{13}}{a_{11}} \\ a_{31} & a_{32} & a_{33} \end{bmatrix} = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ 0 & a_{22} - \frac{a_{21}a_{12}}{a_{11}} & a_{23} - \frac{a_{21}a_{13}}{a_{11}} \\ 0 & a_{32} - \frac{a_{31}a_{12}}{a_{11}} & a_{33} - \frac{a_{31}a_{13}}{a_{11}} \end{bmatrix}$$


---


$$\Downarrow$$

$$\mathbf{E}_{31} \mathbf{E}_{21} \mathbf{A} = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ 0 & \tilde{a}_{22} & \tilde{a}_{23} \\ 0 & \tilde{a}_{32} & \tilde{a}_{33} \end{bmatrix}$$

# Gauss Elimination



3) Zero  $\tilde{a}_{32}$

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -\frac{\tilde{a}_{32}}{\tilde{a}_{22}} & 1 \end{bmatrix} \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ 0 & \tilde{a}_{22} & \tilde{a}_{23} \\ 0 & \tilde{a}_{32} & \tilde{a}_{33} \end{bmatrix} = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ 0 & \tilde{a}_{22} & \tilde{a}_{23} \\ 0 & 0 & \tilde{a}_{33} - \frac{\tilde{a}_{32}\tilde{a}_{23}}{\tilde{a}_{22}} \end{bmatrix} \equiv \mathbf{U}$$

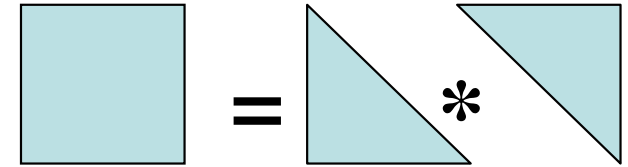
⇓

$$\mathbf{E}_{32} \mathbf{E}_{31} \mathbf{E}_{21} \mathbf{A} = \mathbf{U}$$

⇓

**lower triangular**

# Gauss Elimination



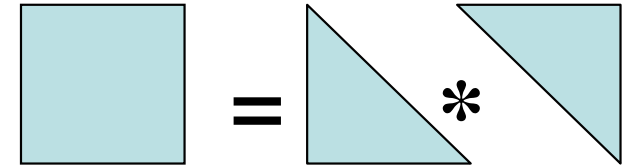
**Claim 1:**

$$\mathbf{E}_{32}\mathbf{E}_{31}\mathbf{E}_{21} = \begin{bmatrix} 1 & 0 & 0 \\ -\frac{a_{21}}{a_{11}} & 1 & 0 \\ -\frac{a_{31}}{a_{11}} & -\frac{\tilde{a}_{32}}{\tilde{a}_{22}} & 1 \end{bmatrix}$$

**Claim 2:**

$$(\mathbf{E}_{32}\mathbf{E}_{31}\mathbf{E}_{21})^{-1} = \begin{bmatrix} 1 & 0 & 0 \\ \frac{a_{21}}{a_{11}} & 1 & 0 \\ \frac{a_{31}}{a_{11}} & \frac{\tilde{a}_{32}}{\tilde{a}_{22}} & 1 \end{bmatrix}$$

# LU Factorization



Therefore, through Gauss Elimination, we have

$$\mathbf{E}_{32} \mathbf{E}_{31} \mathbf{E}_{21} \mathbf{A} = \mathbf{U}$$

$$\mathbf{A} = (\mathbf{E}_{32} \mathbf{E}_{31} \mathbf{E}_{21})^{-1} \mathbf{U}$$

$$\mathbf{A} = \mathbf{L}\mathbf{U}$$

It is called LU factorization. When  $\mathbf{A}$  is symmetric,

$\mathbf{A} = \mathbf{L}\mathbf{L}^T$  which is called *Cholesky Decomposition*

To solve  $Ax=b$  for general form  $A$

To solve  $Ax = b$  becomes

1) Use Gauss elimination to make  $A$  upper triangular, i.e.

$$L^{-1}Ax = L^{-1}b \Rightarrow Ux = z$$

2) Solve  $x$  from  $Ux = z$

This suggests that when doing Gauss elimination, we can do it to the augmented matrix  $[A \ b]$  associated with the linear system.

## An exercise

$$\begin{bmatrix} 2 & -1 & -1 \\ -1 & 2 & -1 \\ -1 & 2 & 2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 4 \end{bmatrix}$$

% build the matrix A

A = [2, -1, 0; -1, 2, -1; 0, -1, 2]

% build the vector b

x\_true = [1:3]';

b = A \* x\_true;

% lu decomposition of A

[l, u] = lu(A)

% solve z from lz = b where z = ux

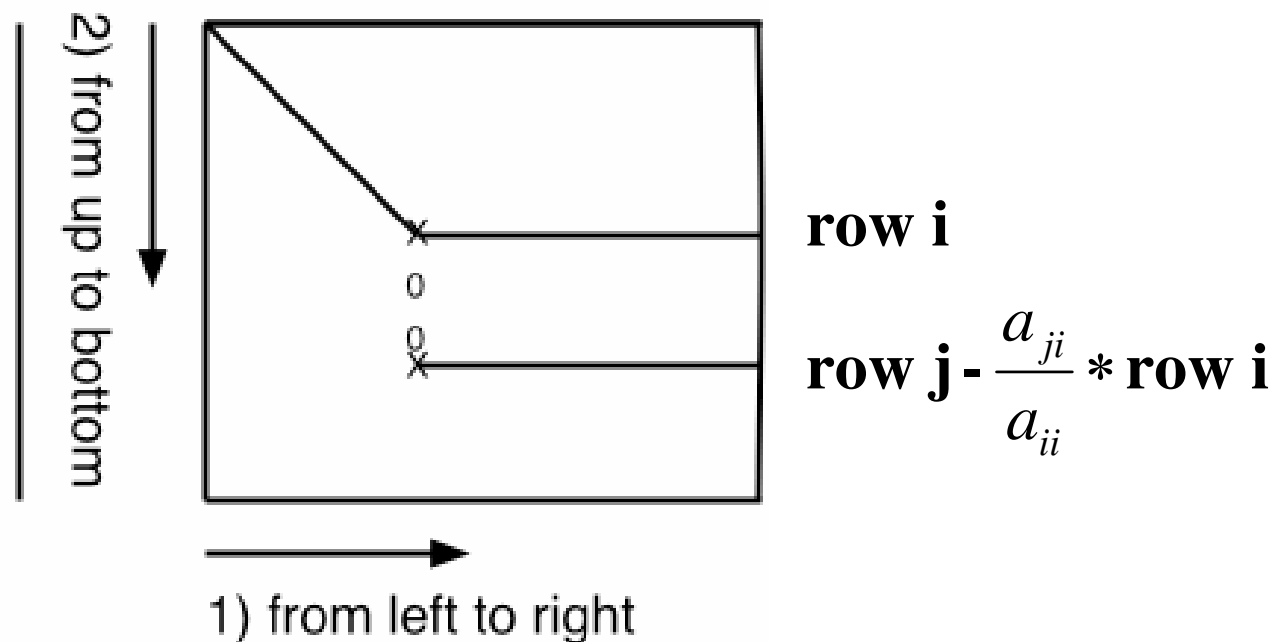
z = l\b;

% solve x from ux = z

x = u\z



# General Gauss Elimination



# What if we have zero or very small diagonal elements?

Sometimes, we need to permute rows of  $A$  so that Gauss elimination can be computed stably, i.e.,  $PA = LU$ . This is called partial pivoting.

For example

$$A = \begin{bmatrix} 0 & 1 \\ 2 & 3 \end{bmatrix}$$

$$PA = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 2 & 3 \end{bmatrix} = \begin{bmatrix} 2 & 3 \\ 0 & 1 \end{bmatrix}$$

$$A = \begin{bmatrix} 0.0001 & 1 \\ 1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 10,000 & 1 \end{bmatrix} \begin{bmatrix} 0.0001 & 1 \\ 0 & -9999 \end{bmatrix} = LU$$

$$PA = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 0.0001 & 1 \\ 1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0.0001 & 1 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 0 & .9999 \end{bmatrix} = LU$$

# Can we always have $A = LU$ ?

No!

If  $\det(\mathbf{A}(1:k,1:k)) \neq 0$  for  $k=1,\dots,n-1$ ,

then  $A \in \mathbb{R}^{n \times n}$  has an LU factorization.

Proof:

See Golub and Loan, Matrix Computation, 3rd edition

# Chapter 5 Finite Difference Methods

## 5.4.1 Crank-Nicolson Method

# Explicit Finite Difference Methods

With  $\frac{\partial f}{\partial t} + rS \frac{\partial f}{\partial S} + \frac{1}{2} \sigma^2 S^2 \frac{\partial^2 f}{\partial S^2} = rf,$

we can obtain an explicit form:

$$\begin{aligned} & \frac{f_{i+1,j} - f_{i,j}}{\Delta t} + O(\Delta t) = \\ & -rj\Delta S \frac{f_{i+1,j+1} - f_{i+1,j-1}}{2\Delta S} - \frac{1}{2} \sigma^2 j^2 (\Delta S)^2 \frac{f_{i+1,j+1} + f_{i+1,j-1} - 2f_{i+1,j}}{\Delta S^2} \\ & + rf_{i+1,j} + O((\Delta S)^2) \end{aligned}$$

# Implicit Finite Difference Methods

With 
$$\frac{\partial f}{\partial t} + rS \frac{\partial f}{\partial S} + \frac{1}{2} \sigma^2 S^2 \frac{\partial^2 f}{\partial S^2} = r f,$$

we can obtain an implicit form :

$$\begin{aligned} \frac{f_{i+1,j} - f_{i,j}}{\Delta t} + O(\Delta t) = \\ -rj\Delta S \frac{f_{i,j+1} - f_{i,j-1}}{2\Delta S} - \frac{1}{2} \sigma^2 j^2 (\Delta S)^2 \frac{f_{i,j+1} + f_{i,j-1} - 2f_{i,j}}{\Delta S^2} \\ + rf_{i,j} + O((\Delta S)^2) \end{aligned}$$

# Crank-Nicolson Methods: average of explicit and implicit methods (Brandmarte, p485)

$$\begin{aligned} \frac{f_{i+1,j} - f_{i,j}}{\Delta t} + O((\Delta t)^2) = & \\ -\frac{rj\Delta S}{2} \left( \frac{f_{i+1,j+1} - f_{i+1,j-1}}{2\Delta S} + \frac{f_{i,j+1} - f_{i,j-1}}{2\Delta S} \right) & \\ -\frac{1}{4}\sigma^2 j^2 (\Delta S)^2 \left( \frac{f_{i+1,j+1} + f_{i+1,j-1} - 2f_{i+1,j}}{\Delta S^2} + \frac{f_{i,j+1} + f_{i,j-1} - 2f_{i,j}}{\Delta S^2} \right) & \\ +\frac{r}{2}(f_{i+1,j} + f_{i,j}) + O((\Delta S)^2) & \end{aligned}$$

## Crank-Nicolson Methods

Rewriting the equation, we get an Crank-Nicolson scheme:

$$\begin{aligned} -\alpha_j f_{i,j-1} + (1 - \beta_j) f_{i,j} - \gamma_j f_{i,j+1} = \\ \alpha_j f_{i+1,j-1} + (1 + \beta_j) f_{i+1,j} + \gamma_j f_{i+1,j+1} \end{aligned} \quad (5.5)$$

where

$$\alpha_j = \frac{\Delta t}{4} (\sigma^2 j^2 - rj)$$

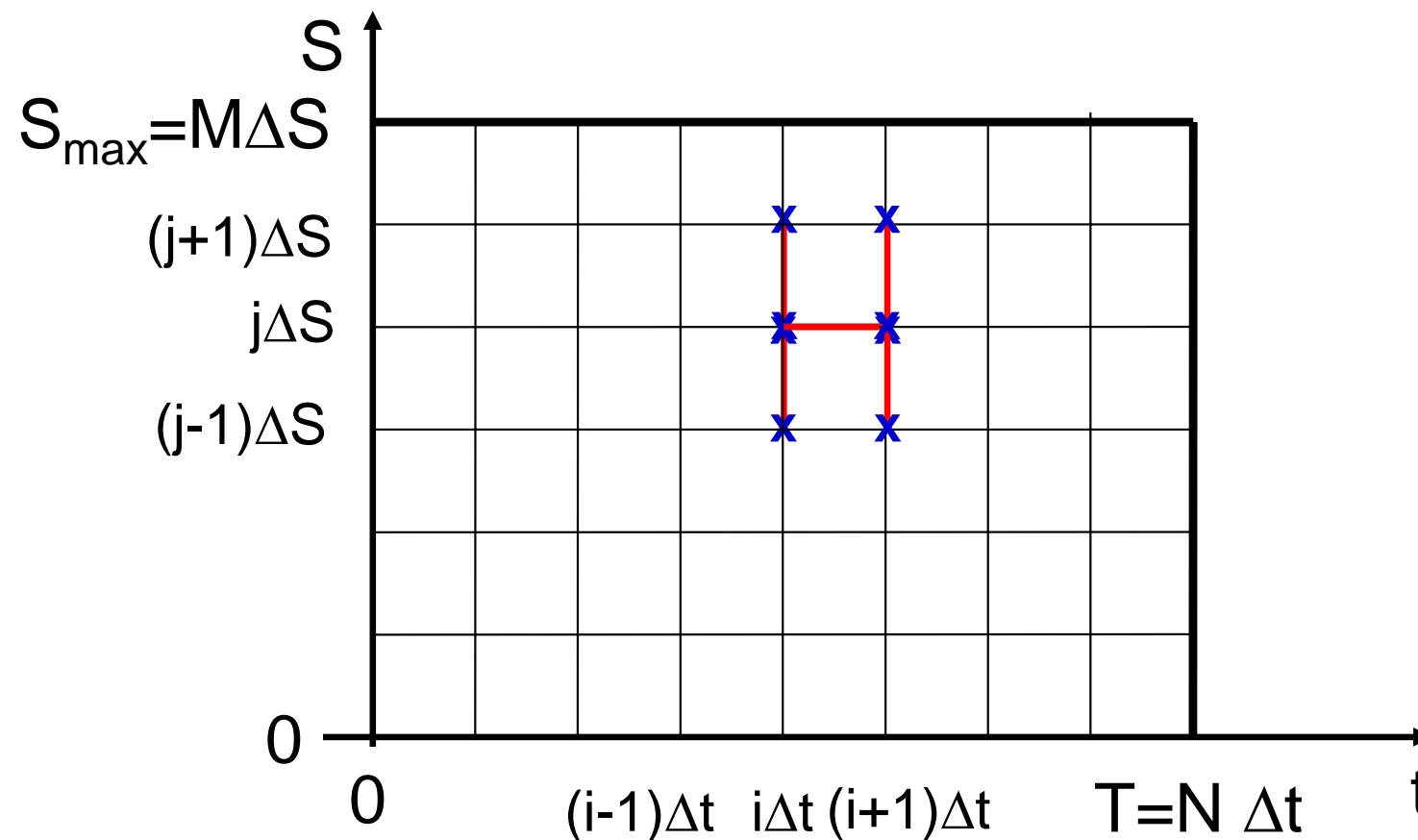
$$\beta_j = -\frac{\Delta t}{2} (\sigma^2 j^2 + r)$$

$$\gamma_j = \frac{\Delta t}{4} (\sigma^2 j^2 + rj)$$

for  $i = N - 1, N - 2, \dots, 1, 0$  and  $j = 1, 2, \dots, M - 1$ .



# Numerical Computation Dependency



# Implementation

Equation (5.5) can be rewritten in matrix form:

$$\mathbf{M}_1 \mathbf{f}_i = \mathbf{M}_2 \mathbf{f}_{i+1} + \mathbf{b}$$

where  $\mathbf{f}_i$  and  $\mathbf{b}_i$  are  $(M-1)$  dimensional vectors

$$\mathbf{f}_i = [f_{i,1}, f_{i,2}, f_{i,3} \cdots, f_{i,M-1}]^T,$$

$$\mathbf{b} = [\alpha_1 (f_{i,0} + f_{i+1,0}), 0 \cdots, \gamma_{M-1} (f_{i,M} + f_{i+1,M})]^T$$

and  $\mathbf{M}_1$  and  $\mathbf{M}_2$  are  $(M-1) \times (M-1)$  symmetric matrices

$$\mathbf{M}_1 = \begin{bmatrix} 1-\beta_1 & -\gamma_1 & 0 & \cdots & 0 \\ -\alpha_2 & 1-\beta_2 & \gamma_2 & \ddots & 0 \\ 0 & \alpha_3 & & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & -\gamma_{M-2} \\ 0 & \cdots & 0 & -\alpha_{M-1} & 1-\beta_{M-1} \end{bmatrix}$$

# Implementation

and

$$\mathbf{M}_2 = \begin{bmatrix} 1 + \beta_1 & \gamma_1 & 0 & \cdots & 0 \\ \alpha_2 & 1 + \beta_2 & \gamma_2 & \ddots & 0 \\ 0 & \alpha_3 & & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \gamma_{M-2} \\ 0 & \cdots & 0 & \alpha_{M-1} & 1 + \beta_{M-1} \end{bmatrix}$$

# Example

We compare Crank-Nicolson Methods for a European put with the exact Black-Scholes formula, where  $T = 5/12$  yr,  $S_0 = \$50$ ,  $K = \$50$ ,  $\sigma = 30\%$ ,  $r = 10\%$ .

Black-Scholes Price: \$2.8446

CN Method with  $S_{\max} = \$100$ ,  $\Delta S = 2$ ,  $\Delta t = 5/1200$ : \$2.8241

CN Method with  $S_{\max} = \$100$ ,  $\Delta S = 1$ ,  $\Delta t = 5/4800$ : \$2.8395

## Example (Stability)

We compare Crank-Nicolson Method for a European put with the exact Black-Scholes formula, where  $T = 5/12$  yr,  $S_0 = \$50$ ,  $K = \$50$ ,  $\sigma = 30\%$ ,  $r = 10\%$ .

Black-Scholes Price: \$2.8446

CN Method with  $S_{\max} = \$100$ ,  $\Delta S = 1.5$ ,  $\Delta t = 5/1200$ : \$3.1370

CN Method with  $S_{\max} = \$100$ ,  $\Delta S = 1$ ,  $\Delta t = 5/1200$ : \$2.8395

# Example: Barrier Options

Barrier options are options where the payoff depends on whether the underlying asset's price reaches a certain level during a certain period of time.

Types:

knock-out: option ceases to exist when the asset price reaches a barrier

Knock-in: option comes into existence when the asset price reaches a barrier

# Example: Barrier Option

We compare Crank-Nicolson Method for a European down-and-out put, where  $T = 5/12$  yr,  $S_0 = \$50$ ,  $K = \$50$ ,  $S_b = \$50$ ,  $\sigma = 40\%$ ,  $r = 10\%$ .

What are the boundary conditions for  $S$ ?

## Example: Barrier Option

We compare Crank-Nicolson Method for a European down-and-out put, where  $T = 5/12$  yr,  $S_0 = \$50$ ,  $K = \$50$ ,  $S_b = \$50$ ,  $\sigma = 40\%$ ,  $r = 10\%$ .

Boundary conditions are  $f(t, S_{max}) = 0$  and  $f(t, S_b) = 0$

Exact Price (Hall, p533-535): \$0.5424

CN Method with  $S_{max} = \$100$ ,  $\Delta S = 0.5$ ,  $\Delta t = 1/1200$ : \$0.5414



# Appendix A.

## Matrix Norms

# Vector Norms

- Norms serve as a way to measure the length of a vector or a matrix
- A vector norm is a function mapping  $\mathbf{x} \in \mathfrak{R}^n$  to a real number  $\|\mathbf{x}\|$  s.t.
  - $\|\mathbf{x}\| > 0$  for any  $\mathbf{x} \neq \mathbf{0}$ ;  $\|\mathbf{x}\| = 0$  iff  $\mathbf{x} = \mathbf{0}$
  - $\|c \mathbf{x}\| = |c| \|\mathbf{x}\|$  for any  $c \in \mathfrak{R}$
  - $\|\mathbf{x} + \mathbf{y}\| \leq \|\mathbf{x}\| + \|\mathbf{y}\|$  for any  $\mathbf{x}, \mathbf{y} \in \mathfrak{R}^n$
- There are various ways to define a norm

$$\|\mathbf{x}\|_p \equiv \left( \sum_{i=1}^n |x_i|^p \right)^{1/p} \quad (p = 2 \text{ is the Euclidean norm})$$

$$\|\mathbf{x}\|_\infty \equiv \max_{1 \leq i \leq n} |x_i|$$

– For example,  $\mathbf{v} = [2 \ 4 \ -1 \ 3]$ .  $\|\mathbf{v}\|_1 = ?$ ,  $\|\mathbf{v}\|_\infty = ?$ ,  $\|\mathbf{v}\|_2 = ?$

# Matrix Norms

- Similarly, a matrix norm is a function mapping  $\mathbf{A} \in \mathfrak{R}^{m \times n}$  to a real number  $\|\mathbf{A}\|$  s.t.

- $\|\mathbf{A}\| > 0$  for any  $\mathbf{A} \neq \mathbf{0}$ ;  $\|\mathbf{A}\| = 0$  iff  $\mathbf{A} = \mathbf{0}$
- $\|c \mathbf{A}\| = |c| \|\mathbf{A}\|$  for any  $c \in \mathfrak{R}$
- $\|\mathbf{A} + \mathbf{B}\| \leq \|\mathbf{A}\| + \|\mathbf{B}\|$  for any  $\mathbf{A}, \mathbf{B} \in \mathfrak{R}^{m \times n}$

- Various commonly used matrix norms

$$\|\mathbf{A}\|_p \equiv \sup_{\mathbf{x} \neq \mathbf{0}} \frac{\|\mathbf{A}\mathbf{x}\|_p}{\|\mathbf{x}\|_p} \qquad \|\mathbf{A}\|_F \equiv \left( \sum_{i=1}^m \sum_{j=1}^n |a_{ij}|^2 \right)^{1/2}$$

$$\|\mathbf{A}\|_1 \equiv \max_{1 \leq j \leq n} \sum_{i=1}^m |a_{ij}| \qquad \|\mathbf{A}\|_\infty \equiv \max_{1 \leq i \leq m} \sum_{j=1}^n |a_{ij}|$$

$$\|\mathbf{A}\|_2 \equiv \sqrt{\rho(\mathbf{A}^T \mathbf{A})}, \text{ the spectral norm, where}$$

$$\rho(\mathbf{B}) \equiv \max \{ |\lambda_k| : \lambda_k \text{ is an eigenvalue of } \mathbf{B} \}$$

## An Example

$$\mathbf{A} = \begin{bmatrix} 2 & 4 & -1 \\ 3 & 1 & 5 \\ -2 & 3 & -1 \end{bmatrix}$$

$$\|\mathbf{A}\|_{\infty} = ? \quad \|\mathbf{A}\|_2 = ?$$

$$\|\mathbf{A}\|_1 = ? \quad \|\mathbf{A}\|_F = ?$$

# Basic Properties of Norms

Let  $\mathbf{A}, \mathbf{B} \in \mathfrak{R}^{n \times n}$  and  $\mathbf{x}, \mathbf{y} \in \mathfrak{R}^n$ . Then

1.  $\|\mathbf{x}\| \geq 0$ ; and  $\|\mathbf{x}\| = 0 \Leftrightarrow \mathbf{x} = \mathbf{0}$
2.  $\|\mathbf{x} + \mathbf{y}\| \leq \|\mathbf{x}\| + \|\mathbf{y}\|$
3.  $\|\alpha \mathbf{x}\| = |\alpha| \|\mathbf{x}\|$  where  $\alpha$  is a real number
4.  $\|\mathbf{A}\mathbf{x}\| \leq \|\mathbf{A}\| \|\mathbf{x}\|$
5.  $\|\mathbf{AB}\| \leq \|\mathbf{A}\| \|\mathbf{B}\|$

# Condition number of a square matrix

All norms in  $\mathbb{R}^n$  ( $\mathbb{R}^{m \times n}$ ) are equivalent. That is, if  $\|\bullet\|_\alpha$  and  $\|\bullet\|_\beta$  are norms on  $\mathbb{R}^n$ , then  $\exists c_1, c_2 > 0$  such that for all  $\mathbf{x} \in \mathbb{R}^n$ , we have

$$c_1 \|\mathbf{x}\|_\alpha \leq \|\mathbf{x}\|_\beta \leq c_2 \|\mathbf{x}\|_\alpha$$

**Condition Number of A Matrix:**  $C \equiv \|\mathbf{A}\| \|\mathbf{A}^{-1}\|$ , where  $\mathbf{A} \in \mathbb{R}^{n \times n}$ .

The condition number gives a measure of how close a matrix is close to singular. The bigger the  $C$ , the harder it is to solve  $\mathbf{Ax} = \mathbf{b}$ .

# Convergence

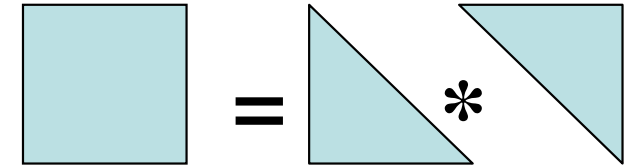
- **vectors  $\mathbf{x}^k$  converges to  $\mathbf{x} \Leftrightarrow \|\mathbf{x}^k - \mathbf{x}\|$  converges to 0**
- **matrix  $\mathbf{A}^k \rightarrow \mathbf{0} \Leftrightarrow \|\mathbf{A}^k - \mathbf{0}\| \rightarrow 0$**

# Appendix B.

## Basic Row Operations



# Basic row operations

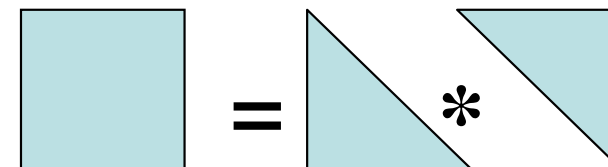


Three kinds of basic row operations:

1) Interchange the order of two rows or (equations)

$$\left[ \begin{array}{ccc|ccc} 0 & 1 & 0 & a_{11} & a_{12} & a_{13} \\ 1 & 0 & 0 & a_{21} & a_{22} & a_{23} \\ 0 & 0 & 1 & a_{31} & a_{32} & a_{33} \end{array} \right] = \left[ \begin{array}{ccc|ccc} a_{21} & a_{22} & a_{23} \\ a_{11} & a_{12} & a_{13} \\ a_{31} & a_{32} & a_{33} \end{array} \right]$$

# Basic row operations



2) Multiply a row by a nonzero constant

$$\left[ \begin{array}{ccc|ccc} c & 0 & 0 & a_{11} & a_{12} & a_{13} \\ 0 & 1 & 0 & a_{21} & a_{22} & a_{23} \\ 0 & 0 & 1 & a_{31} & a_{32} & a_{33} \end{array} \right] = \left[ \begin{array}{ccc} ca_{11} & ca_{12} & ca_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{array} \right]$$

3) Add or subtract rows

$$\left[ \begin{array}{ccc|ccc} 1 & 0 & 0 & a_{11} & a_{12} & a_{13} \\ -1 & 1 & 0 & a_{21} & a_{22} & a_{23} \\ 0 & 0 & 1 & a_{31} & a_{32} & a_{33} \end{array} \right] = \left[ \begin{array}{ccc} a_{11} & a_{12} & a_{13} \\ a_{21} - a_{11} & a_{22} - a_{12} & a_{23} - a_{13} \\ a_{31} & a_{32} & a_{33} \end{array} \right]$$